

# Boolean Logic

## المنطق البولياني

Epp, sections 1.1 and 1.2



# Applications of Boolean logic

- Computer programs
- And computer addition
- Logic problems
- Sudoku

# Boolean propositions

- A proposition is a statement that can be either true or false
  - “The sky is blue”
  - “I is a Engineering major”
  - “ $x == y$ ”
- Not propositions:
  - “Are you Bob?”
  - “ $x := 7$ ”

# Boolean variables

- We use Boolean variables to refer to propositions
  - Usually are lower case letters starting with p (i.e.  $p$ ,  $q$ ,  $r$ ,  $s$ , etc.)
  - A Boolean variable can have one of two values true (T) or false (F)
- A proposition can be...
  - A single variable:  $p$
  - An operation of multiple variables:  $p \wedge (q \vee \neg r)$



# Introduction to Logical Operators

- About a dozen logical operators
  - Similar to algebraic operators + \* - /
- In the following examples,
  - $p$  = “Today is Friday”
  - $q$  = “Today is my birthday”

# Logical operators: Not

- A not operation switches (negates) the truth value
- Symbol:  $\neg$  or  $\sim$
- In C++ and Java, the operand is !
- $\neg p$  = “Today is not Friday”

| $p$ | $\neg p$ |
|-----|----------|
| T   | F        |
| F   | T        |

# Logical operators: And

- An and operation is true if both operands are true
- Symbol:  $\wedge$ 
  - It's like the 'A' in And
- In C++ and Java,  
the operand is `&&`
- $p \wedge q =$  "Today is Friday and  
today is my birthday"

| $p$ | $q$ | $p \wedge q$ |
|-----|-----|--------------|
| T   | T   | T            |
| T   | F   | F            |
| F   | T   | F            |
| F   | F   | F            |

# Logical operators: Or

- An or operation is true if either operands are true
- Symbol:  $\vee$
- In C++ and Java, the operand is `||`
- $p \vee q =$  “Today is Friday or today is my birthday (or possibly both)”

| $p$ | $q$ | $p \vee q$ |
|-----|-----|------------|
| T   | T   | T          |
| T   | F   | T          |
| F   | T   | T          |
| F   | F   | F          |



# Logical operators: Exclusive Or

- An exclusive or operation is true if one of the operands are true, but false if both are true
- Symbol:  $\oplus$
- Often called XOR
- $p \oplus q \equiv (p \vee q) \wedge \neg(p \wedge q)$
- In Java, the operand is  $\wedge$  (but not in C++)
- $p \oplus q =$  “Today is Friday or today is my birthday, but not both”

| $p$ | $q$ | $p \oplus q$ |
|-----|-----|--------------|
| T   | T   | F            |
| T   | F   | T            |
| F   | T   | T            |
| F   | F   | F            |



# Inclusive Or versus Exclusive Or

- Do these sentences mean inclusive or exclusive or?
  - Experience with C++ or Java is required
  - Lunch includes soup or salad
  - To enter the country, you need a passport or a driver's license
  - Publish or perish

# Logical operators: Nand and Nor

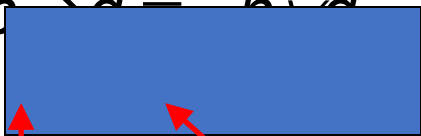
- The negation of And and Or, respectively
- Symbols:  $|$  and  $\downarrow$ , respectively
  - Nand:  $p|q \equiv \neg(p \wedge q)$
  - Nor:  $p \downarrow q \equiv \neg(p \vee q)$

| $p$ | $q$ | $p \wedge q$ | $p \vee q$ | $p q$ | $p \downarrow q$ |
|-----|-----|--------------|------------|-------|------------------|
| T   | T   | T            | T          | F     | F                |
| T   | F   | F            | T          | T     | F                |
| F   | T   | F            | T          | T     | F                |
| F   | F   | F            | F          | T     | T                |

# Logical operators: Conditional 1

- A conditional means “if  $p$  then  $q$ ”
- Symbol:  $\rightarrow$
- $p \rightarrow q$  = “If today is Friday, then today is my birthday”

•  $p \rightarrow q = \neg p \vee q$



the antecedent      the consequence

| $p$ | $q$ | $p \rightarrow q$ | $\neg p \vee q$ |
|-----|-----|-------------------|-----------------|
| T   | T   | T                 | T               |
| T   | F   | F                 | F               |
| F   | T   | T                 | T               |
| F   | F   | T                 | T               |

## Logical operators: Conditional 2

- Let  $p$  = “I am elected” and  $q$  = “I will lower taxes”
- I state:  $p \rightarrow q$  = “If I am elected, then I will lower taxes”
- Consider all possibilities
- Note that if  $p$  is false, then the conditional is true regardless of whether  $q$  is true or false

| $p$ | $q$ | $p \rightarrow q$ |
|-----|-----|-------------------|
| T   | T   | T                 |
| T   | F   | F                 |
| F   | T   | T                 |
| F   | F   | T                 |

# Logical operators: Conditional 3

|     |     |          |          | Conditional       | Inverse                     | Converse          | Contra-<br>positive         |
|-----|-----|----------|----------|-------------------|-----------------------------|-------------------|-----------------------------|
| $p$ | $q$ | $\neg p$ | $\neg q$ | $p \rightarrow q$ | $\neg p \rightarrow \neg q$ | $q \rightarrow p$ | $\neg q \rightarrow \neg p$ |
| T   | T   | F        | F        | T                 | T                           | T                 | T                           |
| T   | F   | F        | T        | F                 | T                           | T                 | F                           |
| F   | T   | T        | F        | T                 | F                           | F                 | T                           |
| F   | F   | T        | T        | T                 | T                           | T                 | T                           |

# Logical operators: Conditional 4

- Alternate ways of stating a conditional:

- $p$  implies  $q$
- If  $p$ ,  $q$
- $p$  is sufficient for  $q$
- $q$  if  $p$
- $q$  whenever  $p$
- $q$  is necessary for  $p$
- $p$  only if  $q$

 I don't like this one

# Logical operators: Bi-conditional 1

- A bi-conditional means “ $p$  if and only if  $q$ ”
- Symbol:  $\leftrightarrow$

• A [redacted]  
 “ [redacted]  
 ( [redacted]

- $p \leftrightarrow q \equiv p \rightarrow q \wedge q \rightarrow p$

- Note that a bi-conditional has the opposite truth values of the exclusive or

| $p$ | $q$ | $p \leftrightarrow q$ |
|-----|-----|-----------------------|
| T   | T   | T                     |
| T   | F   | F                     |
| F   | T   | F                     |
| F   | F   | T                     |



## Logical operators: Bi-conditional 2

- Let  $p$  = “You take this class” and  $q$  = “You get a grade”
- Then  $p \leftrightarrow q$  means  
“You take this class if and only if you get a grade”
- Alternatively, it means “If you take this class, then you get a grade and if you get a grade then you take (took) this class”

| $p$ | $q$ | $p \leftrightarrow q$ |
|-----|-----|-----------------------|
| T   | T   | T                     |
| T   | F   | F                     |
| F   | T   | F                     |
| F   | F   | T                     |



# Boolean operators summary

|     |     | not      | not      | and          | or         | xor          | nand    | nor              | conditional       | bi-conditional        |
|-----|-----|----------|----------|--------------|------------|--------------|---------|------------------|-------------------|-----------------------|
| $p$ | $q$ | $\neg p$ | $\neg q$ | $p \wedge q$ | $p \vee q$ | $p \oplus q$ | $p   q$ | $p \downarrow q$ | $p \rightarrow q$ | $p \leftrightarrow q$ |
| T   | T   | F        | F        | T            | T          | F            | F       | F                | T                 | T                     |
| T   | F   | F        | T        | F            | T          | T            | T       | F                | F                 | F                     |
| F   | T   | T        | F        | F            | T          | T            | T       | F                | T                 | F                     |
| F   | F   | T        | T        | F            | F          | F            | T       | T                | T                 | T                     |

• Learn what they mean, don't just memorize the table!

# Precedence of operators

- Just as in algebra, operators have precedence
  - $4+3*2 = 4+(3*2)$ , not  $(4+3)*2$
- Precedence order (from highest to lowest):  $\neg \wedge \vee \rightarrow \leftrightarrow$ 
  - The first three are the most important
- This means that  $p \vee q \wedge \neg r \rightarrow s \leftrightarrow t$  yields:  $(p \vee (q \wedge (\neg r))) \leftrightarrow (s \rightarrow t)$
- Not is *always* performed before any other operation

# Translating English Sentences

- Problem:
  - $p$  = “It is below freezing”
  - $q$  = “It is snowing”
- It is below freezing and it is snowing
- It is below freezing but not snowing
- It is not below freezing and it is not snowing
- It is either snowing or below freezing (or both)
- If it is below freezing, it is also snowing
- It is either below freezing or it is snowing, but it is not snowing if it is below freezing
- That it is below freezing is necessary and sufficient for it to be snowing

$$p \wedge q$$

$$p \wedge \neg q$$

$$\neg p \wedge \neg q$$

$$p \vee q$$

$$p \rightarrow q$$

$$(p \vee q) \wedge (p \rightarrow \neg q)$$

$$p \leftrightarrow q$$

# Translation Example 1

- Heard on the radio:
  - A study showed that there was a correlation between the more children ate dinners with their families and lower rate of substance abuse by those children
  - Announcer conclusions:
    - If children eat more meals with their family, they will have lower substance abuse
    - If they have a higher substance abuse rate, then they did not eat more meals with their family

# Translation Example 1

- Let  $p$  = “Child eats more meals with family”
- Let  $q$  = “Child has less substance abuse”
- Announcer conclusions:
  - If children eat more meals with their family, they will have lower substance abuse
    - $p \rightarrow q$
  - If they have a higher substance abuse rate, then they did not eat more meals with their family
    - $\neg q \rightarrow \neg p$
- Note that  $p \rightarrow q$  and  $\neg q \rightarrow \neg p$  are logically equivalent

# Translation Example 1

- Let  $p$  = “Child eats more meals with family”
- Let  $q$  = “Child has less substance abuse”
- Remember that the study showed a *correlation*, not a *causation*

| $p$ | $q$ | result | conclusion |
|-----|-----|--------|------------|
| T   | T   | T      | T          |
| T   | F   | ?      | F          |
| F   | T   | ?      | T          |
| F   | F   | T      | T          |

## Translation Example 2

- “I have neither given nor received help on this exam”
  - Rephrased: “I have not given nor received ...”
  - Let  $p$  = “I have given help on this exam”
  - Let  $q$  = “I have received help on this exam”
- Translation is:  $\neg p \downarrow q$

| $p$ | $q$ | $\neg p$ | $\neg p \downarrow q$ |
|-----|-----|----------|-----------------------|
| T   | T   | F        | F                     |
| T   | F   | F        | T                     |
| F   | T   | T        | F                     |
| F   | F   | T        | F                     |



## Translation Example 2

- What they mean is “I have not given and I have not received help on this exam”
  - Or “I have not (given nor received) help on this exam”

| $p$ | $q$ | $\neg p \wedge \neg q$ | $\neg(p \downarrow q)$ |
|-----|-----|------------------------|------------------------|
| T   | T   | F                      | F                      |
| T   | F   | F                      | F                      |
| F   | T   | F                      | F                      |
| F   | F   | T                      | T                      |

- The problem:  $\neg$  has a higher precedence than  $\downarrow$ , but not always in English
- Also, “neither” is vague

# Tautology and Contradiction

- A tautology is a statement that is always true
  - $p \vee \neg p$  will always be true (Negation Law)
- A contradiction is a statement that is always false
  - $p \wedge \neg p$  will always be false (Negation Law)

| $p$ | $p \vee \neg p$ | $p \wedge \neg p$ |
|-----|-----------------|-------------------|
| T   | T               | F                 |
| F   | T               | F                 |

# Logical Equivalence

- A logical equivalence means that the two sides always have the same truth values
  - Symbol is  $\equiv$  or  $\Leftrightarrow$ 
    - We'll use  $\equiv$ , so as not to confuse it with the bi-conditional

# Logical Equivalences of And

- $p \wedge \mathbf{T} \equiv p$

Identity law

| $p$          | $\mathbf{T}$ | $p \wedge \mathbf{T}$ |
|--------------|--------------|-----------------------|
| $\mathbf{T}$ | $\mathbf{T}$ | $\mathbf{T}$          |
| $\mathbf{F}$ | $\mathbf{T}$ | $\mathbf{F}$          |

- $p \wedge \mathbf{F} \equiv \mathbf{F}$

Domination law

| $p$          | $\mathbf{F}$ | $p \wedge \mathbf{F}$ |
|--------------|--------------|-----------------------|
| $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$          |
| $\mathbf{F}$ | $\mathbf{F}$ | $\mathbf{F}$          |

# Logical Equivalences of And

- $p \wedge p \equiv p$

Idempotent law

| $p$ | $p$ | $p \wedge p$ |
|-----|-----|--------------|
| T   | T   | T            |
| F   | F   | F            |

- $p \wedge q \equiv q \wedge p$

Commutative law

| $p$ | $q$ | $p \wedge q$ | $q \wedge p$ |
|-----|-----|--------------|--------------|
| T   | T   | T            | T            |
| T   | F   | F            | F            |
| F   | T   | F            | F            |
| F   | F   | F            | F            |

# Logical Equivalences of And

- $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$       Associative law

| p | q | r | $p \wedge q$ | $(p \wedge q) \wedge r$ | $q \wedge r$ | $p \wedge (q \wedge r)$ |
|---|---|---|--------------|-------------------------|--------------|-------------------------|
| T | T | T | T            | T                       | T            | T                       |
| T | T | F | T            | F                       | F            | F                       |
| T | F | T | F            | F                       | F            | F                       |
| T | F | F | F            | F                       | F            | F                       |
| F | T | T | F            | F                       | T            | F                       |
| F | T | F | F            | F                       | F            | F                       |
| F | F | T | F            | F                       | F            | F                       |
| F | F | F | F            | F                       | F            | F                       |

# Logical Equivalences of Or

- $p \vee \mathbf{T} \equiv \mathbf{T}$

- $p \vee \mathbf{F} \equiv p$

- $p \vee p \equiv p$

- $p \vee q \equiv q \vee p$

- $(p \vee q) \vee r \equiv p \vee (q \vee r)$

Identity law

Domination law

Idempotent law

Commutative law

Associative law



# Corollary of the Associative Law

- $(p \wedge q) \wedge r \equiv p \wedge q \wedge r$
- $(p \vee q) \vee r \equiv p \vee q \vee r$
- Similar to  $(3+4)+5 = 3+4+5$
- Only works if ALL the operators are the same!





# Logical Equivalences of Not

- $\neg(\neg p) \equiv p$

Double negation law

- $p \vee \neg p \equiv T$

Negation law

- $p \wedge \neg p \equiv F$

Negation law

# DeMorgan's Law

- Probably the most important logical equivalence
- To negate  $p \wedge q$  (or  $p \vee q$ ), you “flip” the sign, and negate BOTH  $p$  and  $q$ 
  - Thus,  $\neg(p \wedge q) \equiv \neg p \vee \neg q$
  - Thus,  $\neg(p \vee q) \equiv \neg p \wedge \neg q$

| p | q | $\neg p$ | $\neg q$ | $p \wedge q$ | $\neg(p \wedge q)$ | $\neg p \vee \neg q$ | $p \vee q$ | $\neg(p \vee q)$ | $\neg p \wedge \neg q$ |
|---|---|----------|----------|--------------|--------------------|----------------------|------------|------------------|------------------------|
| T | T | F        | F        | T            | F                  | F                    | T          | F                | F                      |
| T | F | F        | T        | F            | T                  | T                    | T          | F                | F                      |
| F | T | T        | F        | F            | T                  | T                    | T          | F                | F                      |
| F | F | T        | T        | F            | T                  | T                    | F          | T                | T                      |

# Yet more equivalences

- Distributive:

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

- Absorption

$$p \vee (p \wedge q) \equiv p$$

$$p \wedge (p \vee q) \equiv p$$

# How to prove two propositions are equivalent?

- Two methods:
  - Using truth tables
    - Not good for long formulae
    - In this course, only allowed if specifically stated!
  - Using the logical equivalences
    - The preferred method
- Example: show that:

$$(p \rightarrow r) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$$

# Using Truth Tables

$$(p \rightarrow r) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$$

| p | q | r | $p \rightarrow r$ | $q \rightarrow r$ | $(p \rightarrow r) \vee (q \rightarrow r)$ | $p \wedge q$ | $(p \wedge q) \rightarrow r$ |
|---|---|---|-------------------|-------------------|--|--------------|------------------------------|
| T | T | T | T                 | T                 |  | T            |                              |
| T | T | F | F                 | F                 |  | T            |                              |
| T | F | T | T                 | T                 |  | F            |                              |
| T | F | F | F                 | T                 |  | F            |                              |
| F | T | T | T                 | T                 |  | F            |                              |
| F | T | F | T                 | F                 |  | F            |                              |
| F | F | T | T                 | T                 |  | F            |                              |
| F | F | F | T                 | T                 |  | F            |                              |

# Using Logical Equivalences

$$\underline{(p \rightarrow r)} \vee \underline{(q \rightarrow r)} \equiv \underline{(p \wedge q) \rightarrow r} \quad \text{Original statement}$$

$$\underline{(\neg p \vee r)} \wedge \underline{(\neg q \vee r)} \equiv \underline{(p \wedge q) \rightarrow r} \equiv \neg p \vee q$$

Definition of implication

DeMorgan's Law  $\neg(p \wedge q) \equiv \neg p \vee \neg q$

$$\underline{\neg p \vee r} \vee \underline{\neg q \vee r} \equiv \underline{(\neg p \vee r) \vee (\neg q \vee r)} \equiv \neg p \vee r \vee \neg q \vee r$$

Associativity of  $\vee$

$$\underline{\neg p \vee r} \vee \underline{\neg q \vee r} \equiv \neg p \vee \neg q \vee r$$

Re-arranging

$$\underline{\neg p \vee r} \vee \underline{\neg q \vee r} \equiv \underline{\neg p \vee \neg q} \vee r$$

Idempotent Law

# Logical Thinking

- At a trial:
  - Bill says: “Sue is guilty and Fred is innocent.”
  - Sue says: “If Bill is guilty, then so is Fred.”
  - Fred says: “I am innocent, but at least one of the others is guilty.”
- Let  $b$  = Bill is innocent,  $f$  = Fred is innocent, and  $s$  = Sue is innocent
- Statements are:
  - $\neg s \wedge f$
  - $\neg b \rightarrow \neg f$
  - $f \wedge (\neg b \vee \neg s)$

# Can all of their statements be true?

- Show:  $(\neg s \wedge f) \wedge (\neg b \rightarrow \neg f) \wedge (f \wedge (\neg b \vee \neg s))$

| b | f | s | $\neg b$ | $\neg f$ | $\neg s$ | $\neg s \wedge f$ | $\neg b \rightarrow \neg f$ |
|---|---|---|----------|----------|----------|-------------------|-----------------------------|
| T | T | T | F        | F        | F        | F                 | T                           |
|   |   |   |          |          |          |                   |                             |
| T | F | T | F        | T        | F        | F                 | T                           |
| T | F | F | F        | T        | T        | F                 | T                           |
| F | T | T | T        | F        | F        | F                 | F                           |
| F | T | F | T        | F        | T        | T                 | F                           |
| F | F | T | T        | T        | F        | F                 | T                           |
| F | F | F | T        | T        | T        | F                 | T                           |

| $f \wedge (\neg b \vee \neg s)$ |
|---------------------------------|
| F                               |
|                                 |
| F                               |
| F                               |
| T                               |
| T                               |
| F                               |
| F                               |



Are all of their statements true?

Show values for  $s$ ,  $b$ , and  $f$  such that the equation is true

|  |                           |
|--|---------------------------|
| $(\neg s \wedge f) \wedge (\neg b \rightarrow \neg f) \wedge (f \wedge (\neg b \vee \neg s)) \equiv T$ | Original statement        |
| $(\neg s \wedge f) \wedge (b \vee \neg f) \wedge (f \wedge (\neg b \vee \neg s)) \equiv T$             | Definition of implication |
| $\neg s \wedge f \wedge (b \vee \neg f) \wedge f \wedge (\neg b \vee \neg s) \equiv T$                 | Associativity of AND      |
| $\neg s \wedge f \wedge f \wedge (b \vee \neg f) \wedge (\neg b \vee \neg s) \equiv T$                 | Re-arranging              |
| $\neg s \wedge f \wedge (b \vee \neg f) \wedge (\neg b \vee \neg s) \equiv T$                          | Idempotent law            |
| $f \wedge (b \vee \neg f) \wedge \neg s \wedge (\neg s \vee \neg b) \equiv T$                          | Re-arranging              |
| $f \wedge (b \vee \neg f) \wedge \neg s \equiv T$  | Absorption law            |
| $(f \wedge (b \vee \neg f)) \wedge \neg s \equiv T$  | Re-arranging              |
| $((f \wedge b) \vee (f \wedge \neg f)) \wedge \neg s \equiv T$   | Distributive law          |
| $((f \wedge b) \vee F) \wedge \neg s \equiv T$   | Negation law              |
| $(f \wedge b) \wedge \neg s \equiv T$  | Domination law            |
| $f \wedge b \wedge \neg s \equiv T$  | Associativity of AND      |

What if it weren't possible to assign such values to  $s$ ,  $b$ , and  $f$ ?

$$(\neg s \wedge f) \wedge (\neg b \rightarrow \neg f) \wedge (f \wedge (\neg b \vee \neg s)) \wedge s = T$$

Original statement

$$(\neg s \wedge f) \wedge (b \vee \neg f) \wedge (f \wedge (\neg b \vee \neg s)) \wedge s = T$$

Definition of implication

... (same as previous slide)

$$(f \wedge b) \wedge \neg s \wedge s = T$$

Domination law

$$f \wedge b \wedge \neg s \wedge s = T$$

Re-arranging

$$f \wedge b \wedge F = T$$

Negation law

$$f \wedge F = T$$

Domination law

$$F = T$$

Domination law

Contradiction!

# Functional completeness

- All the “extended” operators have equivalences using only the 3 basic operators (and, or, not)
  - The extended operators: nand, nor, xor, conditional, bi-conditional
- Given a limited set of operators, can you write an equivalence of the 3 basic operators?
  - If so, then that group of operators is functionally complete

coffee "or" tea ← ? exclusive-or

How to construct a compound statement for exclusive-or?

| p | q | p <span style="border: 1px solid black; padding: 2px;">?</span> q |
|---|---|---|
| T | T | F   |
| T | F | T   |
| F | T | T   |
| F | F | F   |

Idea 1: Look at the true rows

$$(p \wedge \neg q) \vee (\neg p \wedge q)$$

Idea 2: Look at the false rows

$$\neg(p \wedge q) \wedge \neg(\neg p \wedge \neg q)$$

Idea 3: Guess and check

$$(p \vee q) \wedge \neg(p \wedge q)$$

# Logical Equivalence



$$p \oplus q \equiv (p \vee q) \wedge \neg(p \wedge q)$$

| $p$ | $q$ | $p \oplus q$ | $p \vee q$ | $\neg(p \wedge q)$ |   |
|-----|-----|--------------|------------|--------------------|---|
| T   | T   | F            | T          | F                  | F |
| T   | F   | T            | T          | T                  | T |
| F   | T   | T            | T          | T                  | T |
| F   | F   | F            | F          | T                  | F |

**Logical equivalence:** Two statements have the same truth table

# Writing Logical Formula for a Truth Table



Given a truth table, how to write a logical formula with the same function?

First write down a small formula for each row, so that the formula is true if the inputs are exactly the same as the row.

Then use idea 1 or idea 2.

| p | q | r | output |
|---|---|---|--------|
| T | T | T | F      |
| T | T | F | T      |
| T | F | T | T      |
| T | F | F | F      |
| F | T | T | T      |
| F | T | F | T      |
| F | F | T | T      |
| F | F | F | F      |

$$p \wedge q \wedge r$$

$$p \wedge q \wedge \neg r$$

$$p \wedge \neg q \wedge r$$

$$p \wedge \neg q \wedge \neg r$$

$$\vee(\neg p \wedge \neg q \wedge r)$$

$$\neg p \wedge q \wedge \neg r$$

$$\neg p \wedge \neg q \wedge r$$

$$\neg p \wedge \neg q \wedge \neg r$$

Idea 1: Look at the true rows and take the “or”.

$$(p \wedge q \wedge \neg r)$$

$$\vee(p \wedge \neg q \wedge r)$$

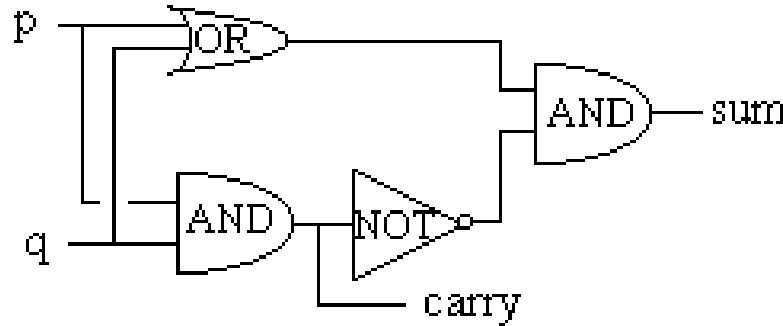
$$\vee(\neg p \wedge q \wedge r)$$

$$\vee(\neg p \wedge q \wedge \neg r)$$

The formula is true iff the input is one of the true rows.

# Writing Logical Formula for a Truth Table

Digital logic:



| p | q | sum | carry |
|---|---|-----|-------|
| 1 | 1 | 0   | 1     |
| 1 | 0 | 1   | 0     |
| 0 | 1 | 1   | 0     |
| 0 | 0 | 0   | 0     |

Idea 2: Look at the false rows, **negate** and take the **“and”**.

| p | q | r | output |
|---|---|---|--------|
| T | T | T | F      |
| T | T | F | T      |
| T | F | T | T      |
| T | F | F | F      |
| F | T | T | T      |
| F | T | F | T      |
| F | F | T | T      |
| F | F | F | F      |

- $p \wedge q \wedge r$
- $p \wedge q \wedge \neg r$
- $p \wedge \neg q \wedge r$
- $p \wedge \neg q \wedge \neg r$
- $\neg p \wedge q \wedge r$
- $\neg p \wedge q \wedge \neg r$
- $\neg p \wedge \neg q \wedge r$
- $\neg p \wedge \neg q \wedge \neg r$

$$\neg(p \wedge q \wedge r)$$

$$\wedge \neg(p \wedge \neg q \wedge \neg r)$$

$$\wedge \neg(\neg p \wedge \neg q \wedge \neg r)$$

can be simplified further

The formula is true iff the input is **not** one of the false row.